

Runir's Description Logic Feature Syntax and Semantics

Dominik Drexler
dominik.drexler@liu.se

1 Preliminaries

Let \mathcal{I} be an interpretation with non-empty domain $\Delta^{\mathcal{I}}$. Each atomic concept A is interpreted as a set $A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$, each atomic role R as a binary relation $R^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$, and each individual name a as an element $a^{\mathcal{I}} \in \Delta^{\mathcal{I}}$.

Concept descriptions denote subsets of $\Delta^{\mathcal{I}}$. Role descriptions denote binary relations over $\Delta^{\mathcal{I}}$. For planning states, let s denote the set of atoms true in the current state, and let γ denote the set of goal literals.

2 Concept Constructors

Name	Syntax	Semantics
Top	\top	$\top^{\mathcal{I}} = \Delta^{\mathcal{I}}$
Bottom	\perp	$\perp^{\mathcal{I}} = \emptyset$
Positive atomic state concept	P	$P^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid P(a) \in s\}$
Negative atomic state concept	$\neg P$	$\neg P^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus P^{\mathcal{I}}$
Positive atomic goal concept	G^+	$G^{+\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid G(a) \in \gamma\}$
Negative atomic goal concept	G^-	$G^{-\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \neg G(a) \in \gamma\}$
Intersection	$C \sqcap D$	$C \sqcap D^{\mathcal{I}} = C^{\mathcal{I}} \cap D^{\mathcal{I}}$
Union	$C \sqcup D$	$C \sqcup D^{\mathcal{I}} = C^{\mathcal{I}} \cup D^{\mathcal{I}}$
Negation	$\neg C$	$\neg C^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus C^{\mathcal{I}}$
Value restriction	$\forall R.C$	$\forall R.C^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \forall b. (a, b) \in R^{\mathcal{I}} \Rightarrow b \in C^{\mathcal{I}}\}$
Existential quantification	$\exists R.C$	$\exists R.C^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \exists b. (a, b) \in R^{\mathcal{I}} \wedge b \in C^{\mathcal{I}}\}$
At-least number restriction	$\geq n R$	$\geq n R^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}}\} \geq n\}$
At-most number restriction	$\leq n R$	$\leq n R^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}}\} \leq n\}$
Exact number restriction	$= n R$	$= n R^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}}\} = n\}$
Qualified at-least number restriction	$\geq n R.C$	$\geq n R.C^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}} \wedge b \in C^{\mathcal{I}}\} \geq n\}$
Qualified at-most number restriction	$\leq n R.C$	$\leq n R.C^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}} \wedge b \in C^{\mathcal{I}}\} \leq n\}$

Name	Syntax	Semantics
Qualified exact number restriction	$= n R.C$	$= n R.C^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}} \wedge b \in C^{\mathcal{I}}\} = n\}$
Role-value map inclusion	$R \subseteq S$	$R \subseteq S^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \forall b. (a, b) \in R^{\mathcal{I}} \Rightarrow (a, b) \in S^{\mathcal{I}}\}$
Agreement	$R \doteq S$	$R \doteq S^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \forall b. (a, b) \in R^{\mathcal{I}} \Leftrightarrow (a, b) \in S^{\mathcal{I}}\}$
Role fillers	$R : \{a_1, \dots, a_k\}$	$R : \{a_1, \dots, a_k\}^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} \mid \{a_1^{\mathcal{I}}, \dots, a_k^{\mathcal{I}}\} \subseteq \{b \in \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}}\}\}$
One-of	$\{a_1, \dots, a_k\}$	$\{a_1, \dots, a_k\}^{\mathcal{I}} = \{a_1^{\mathcal{I}}, \dots, a_k^{\mathcal{I}}\}$
Nominal	$\{a\}$	$\{a\}^{\mathcal{I}} = \{a^{\mathcal{I}}\}$

Here $n \in \mathbb{N}$. In role fillers and one-of constructors, the listed object names are interpreted as individual names in \mathcal{I} .

3 Concrete Syntax for Concept Constructors

Name	Concrete syntax	Abstract syntax
Top	<code>c_top</code>	\top
Bottom	<code>c_bot</code>	\perp
Positive atomic state concept	<code>(c_atomic_state "p" true)</code>	P
Negative atomic state concept	<code>(c_atomic_state "p" false)</code>	$\neg P$
Positive atomic goal concept	<code>(c_atomic_goal "p" true)</code>	G^+
Negative atomic goal concept	<code>(c_atomic_goal "p" false)</code>	G^-
Intersection	<code>(c_and C1 ... Cn)</code>	$C_1 \sqcap \dots \sqcap C_n$
Union	<code>(c_or C1 ... Cn)</code>	$C_1 \sqcup \dots \sqcup C_n$
Negation	<code>(c_not C)</code>	$\neg C$
Value restriction	<code>(c_all R C)</code>	$\forall R.C$
Existential quantification	<code>(c_some R C)</code>	$\exists R.C$
At-least number restriction	<code>(c_at_least n R)</code>	$\geq n R$
At-most number restriction	<code>(c_at_most n R)</code>	$\leq n R$
Exact number restriction	<code>(c_exactly n R)</code>	$= n R$
Qualified at-least restriction	<code>(c_at_least n R C)</code>	$\geq n R.C$
Qualified at-most restriction	<code>(c_at_most n R C)</code>	$\leq n R.C$
Qualified exact restriction	<code>(c_exactly n R C)</code>	$= n R.C$
Same-as / agreement	<code>(c_same_as R1 R2)</code>	$R_1 \doteq R_2$
Role-value map	<code>(c_subset R1 R2)</code>	$R_1 \subseteq R_2$
Role fillers	<code>(c_fillers R a1 ... an)</code>	$R : \{a_1, \dots, a_n\}$
One-of	<code>(c_one_of a1 ... an)</code>	$\{a_1, \dots, a_n\}$

Name	Concrete syntax	Abstract syntax
Nominal	(c_nominal a)	{a}

4 Restrictions on Role Interpretations

Functional roles. A feature f is interpreted as a functional binary relation:

$$\forall a, b, c. (a, b) \in f^{\mathcal{I}} \wedge (a, c) \in f^{\mathcal{I}} \Rightarrow b = c.$$

Equivalently, features may be viewed as partial functions.

Transitive roles. A transitive role R is interpreted as a transitive binary relation:

$$\forall a, b, c. (a, b) \in R^{\mathcal{I}} \wedge (b, c) \in R^{\mathcal{I}} \Rightarrow (a, c) \in R^{\mathcal{I}}.$$

5 Role Constructors

Name	Syntax	Semantics
Universal role	U	$U^{\mathcal{I}} = \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$
Positive atomic state role	P	$P^{\mathcal{I}} = \{(a, b) \in \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \mid P(a, b) \in s\}$
Negative atomic state role	$\neg P$	$\neg P^{\mathcal{I}} = (\Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}) \setminus P^{\mathcal{I}}$
Positive atomic goal role	G^+	$G^{+\mathcal{I}} = \{(a, b) \in \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \mid G(a, b) \in \gamma\}$
Negative atomic goal role	G^-	$G^{-\mathcal{I}} = \{(a, b) \in \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \mid \neg G(a, b) \in \gamma\}$
Intersection	$R \sqcap S$	$R \sqcap S^{\mathcal{I}} = R^{\mathcal{I}} \cap S^{\mathcal{I}}$
Union	$R \sqcup S$	$R \sqcup S^{\mathcal{I}} = R^{\mathcal{I}} \cup S^{\mathcal{I}}$
Complement	$\neg R$	$\neg R^{\mathcal{I}} = (\Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}) \setminus R^{\mathcal{I}}$
Inverse	R^-	$(R^-)^{\mathcal{I}} = \{(b, a) \in \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \mid (a, b) \in R^{\mathcal{I}}\}$
Composition	$R \circ S$	$R \circ S^{\mathcal{I}} = \{(a, c) \mid \exists b. (a, b) \in R^{\mathcal{I}} \wedge (b, c) \in S^{\mathcal{I}}\}$
Transitive closure	R^+	$(R^+)^{\mathcal{I}} = \bigcup_{n \geq 1} (R^{\mathcal{I}})^n$
Reflexive-transitive closure	R^*	$(R^*)^{\mathcal{I}} = \bigcup_{n \geq 0} (R^{\mathcal{I}})^n$
Role restriction	$R C$	$R C^{\mathcal{I}} = R^{\mathcal{I}} \cap (\Delta^{\mathcal{I}} \times C^{\mathcal{I}})$
Identity	$\text{id}(C)$	$\text{id}(C)^{\mathcal{I}} = \{(d, d) \mid d \in C^{\mathcal{I}}\}$

The iterated composition $(R^{\mathcal{I}})^n$ is defined by

$$(R^{\mathcal{I}})^0 = \{(d, d) \mid d \in \Delta^{\mathcal{I}}\}, \quad (R^{\mathcal{I}})^{n+1} = (R^{\mathcal{I}})^n \circ R^{\mathcal{I}}.$$

6 Concrete Syntax for Role Constructors

Name	Concrete syntax	Abstract syntax
Universal role	<code>r_universal</code>	U
Positive atomic state role	<code>(r_atomic_state "p" true)</code>	P
Negative atomic state role	<code>(r_atomic_state "p" false)</code>	$\neg P$
Positive atomic goal role	<code>(r_atomic_goal "p" true)</code>	G^+
Negative atomic goal role	<code>(r_atomic_goal "p" false)</code>	G^-
Intersection	<code>(r_and R1 ... Rn)</code>	$R_1 \sqcap \dots \sqcap R_n$
Union	<code>(r_or R1 ... Rn)</code>	$R_1 \sqcup \dots \sqcup R_n$
Complement	<code>(r_complement R)</code>	$\neg R$
Inverse	<code>(r_inverse R)</code>	R^-
Composition	<code>(r_composition R1 ... Rn)</code>	$R_1 \circ \dots \circ R_n$
Transitive closure	<code>(r_transitive_closure R)</code>	R^+
Reflexive-transitive closure	<code>(r_reflexive_transitive_closure R)</code>	R^*
Role restriction	<code>(r_restriction R C)</code>	$R C$
Identity	<code>(r_identity C)</code>	$\text{id}(C)$

7 Boolean Constructors

Name	Syntax	Semantics
Positive atomic state predicate	<code>holds(p)</code>	true iff there exists an atom $p() \in s$
Negative atomic state predicate	<code>\negholds(p)</code>	true iff no current atom with predicate p is true in the state
Positive atomic goal predicate	<code>goal(p)</code>	true iff $p() \in \gamma$
Negative atomic goal predicate	<code>\neggoal(p)</code>	true iff $\neg p() \in \gamma$
Non-empty concept	$C \neq \emptyset$	true iff $C^{\mathcal{I}} \neq \emptyset$
Non-empty role	$R \neq \emptyset$	true iff $R^{\mathcal{I}} \neq \emptyset$

8 Concrete Syntax for Boolean Constructors

Name	Concrete syntax	Abstract syntax
Positive atomic state predicate	(b_atomic_state "p" true)	holds(p)
Negative atomic state predicate	(b_atomic_state "p" false)	\neg holds(p)
Positive atomic goal predicate	(b_atomic_goal "p" true)	goal(p)
Negative atomic goal predicate	(b_atomic_goal "p" false)	\neg goal(p)
Non-empty concept	(b_nonempty C)	$C \neq \emptyset$
Non-empty role	(b_nonempty R)	$R \neq \emptyset$

9 Numerical Constructors

Name	Syntax	Semantics
Concept count	$ C $	$ C^{\mathcal{I}} $
Role count	$ R $	$ R^{\mathcal{I}} $
Role distance	$d_R(C, D)$	$\min\{n \in \mathbb{N} \mid C^{\mathcal{I}} \times D^{\mathcal{I}} \cap (R^{\mathcal{I}})^n \neq \emptyset\}$, or ∞ if empty

10 Concrete Syntax for Numerical Constructors

Name	Concrete syntax	Abstract syntax
Concept count	(n_count C)	$ C $
Role count	(n_count R)	$ R $
Role distance	(n_distance C R D)	$d_R(C, D)$

Source

This document summarizes Appendix 1, Section A1.2 of *The Description Logic Handbook*, edited by Franz Baader, Deborah L. McGuinness, Daniele Nardi, and Peter F. Patel-Schneider.